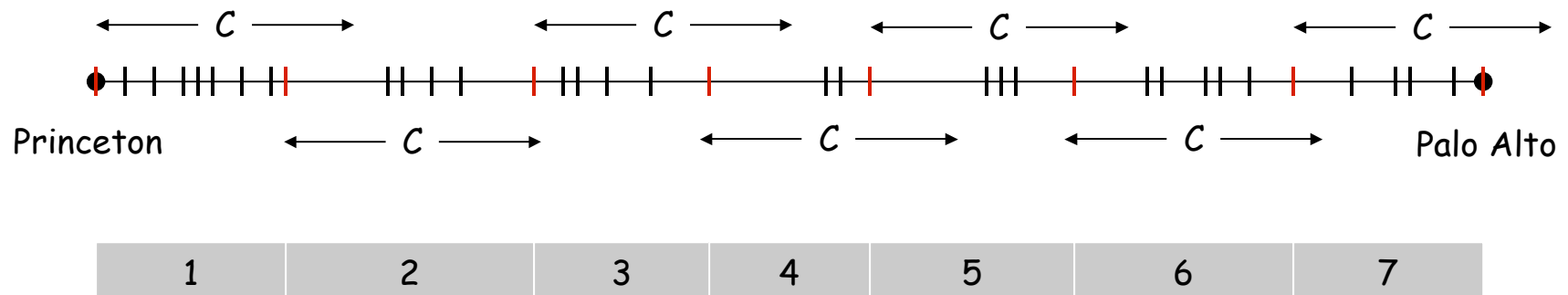


Selecting Breakpoints

Selecting breakpoints.

- Road trip from Princeton to Palo Alto along fixed route.
- Refueling stations at certain points along the way.
- Fuel capacity = C .
- Goal: makes as few refueling stops as possible.

Greedy algorithm. Go as far as you can before refueling.



Selecting Breakpoints: Greedy Algorithm

Truck driver's algorithm.

```
Sort breakpoints so that:  $0 = b_0 < b_1 < b_2 < \dots < b_n = L$ 
```

```
 $S \leftarrow \{0\}$   $\leftarrow$  breakpoints selected
```

```
 $x \leftarrow 0$   $\leftarrow$  current location
```

```
while ( $x \neq b_n$ )
```

```
    let  $p$  be largest integer such that  $b_p \leq x + C$ 
```

```
    if ( $b_p = x$ )
```

```
        return "no solution"
```

```
     $x \leftarrow b_p$ 
```

```
     $S \leftarrow S \cup \{p\}$ 
```

```
return  $S$ 
```

Implementation. $O(n \log n)$

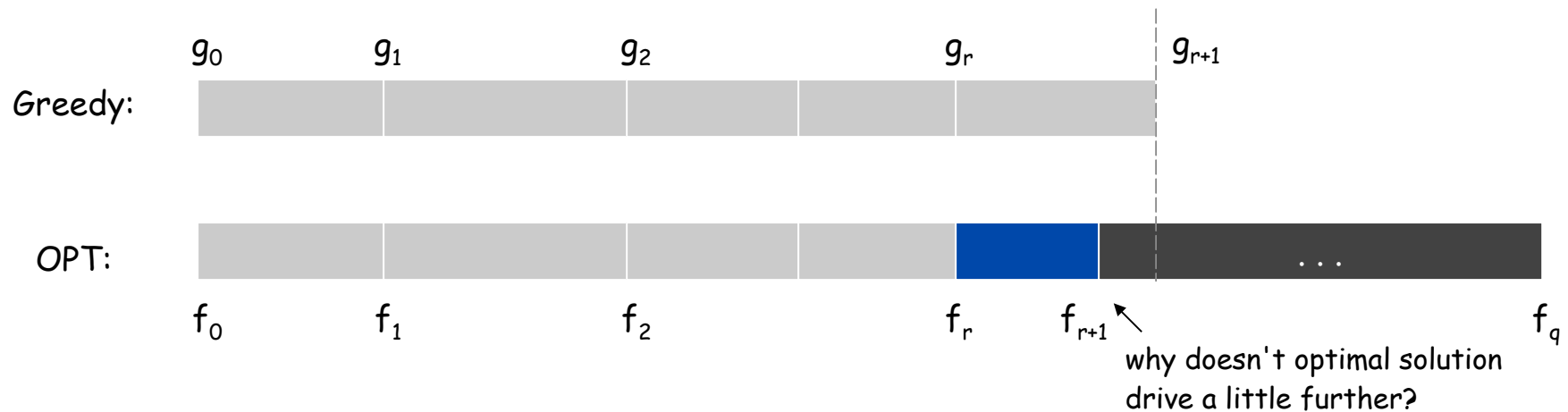
- Use binary search to select each breakpoint p .

Selecting Breakpoints: Correctness

Theorem. Greedy algorithm is optimal.

Pf. (by contradiction)

- Assume greedy is not optimal, and let's see what happens.
- Let $0 = g_0 < g_1 < \dots < g_p = L$ denote set of breakpoints chosen by greedy.
- Let $0 = f_0 < f_1 < \dots < f_q = L$ denote set of breakpoints in an optimal solution with $f_0 = g_0, f_1 = g_1, \dots, f_r = g_r$ for largest possible value of r .
- Note: $g_{r+1} > f_{r+1}$ by greedy choice of algorithm.



Selecting Breakpoints: Correctness

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